COMMUNITY
THE ASF CONFERENCE
CODE

Impalas living on Iceberg

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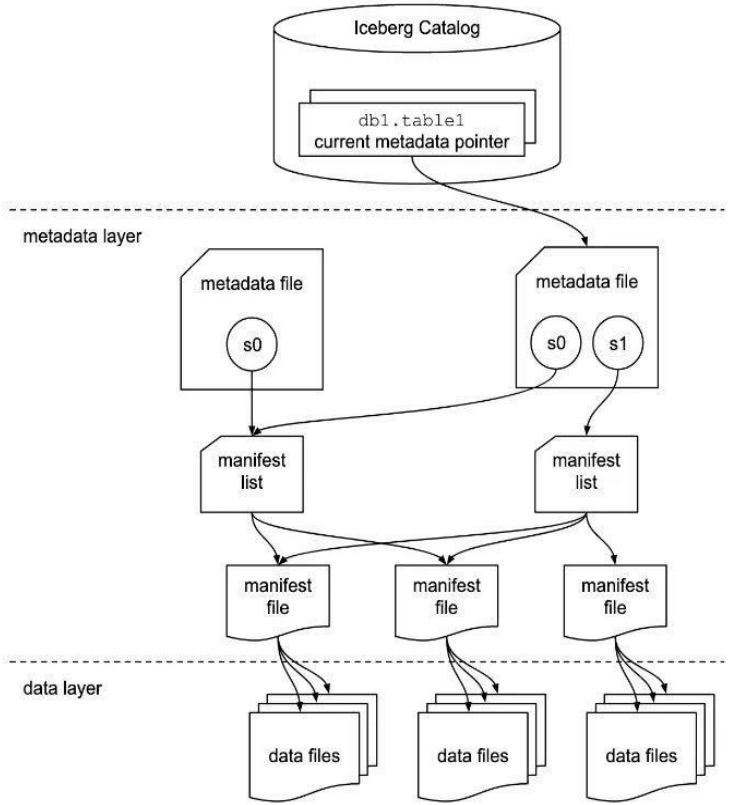
Contents

1. Introduction - Iceberg

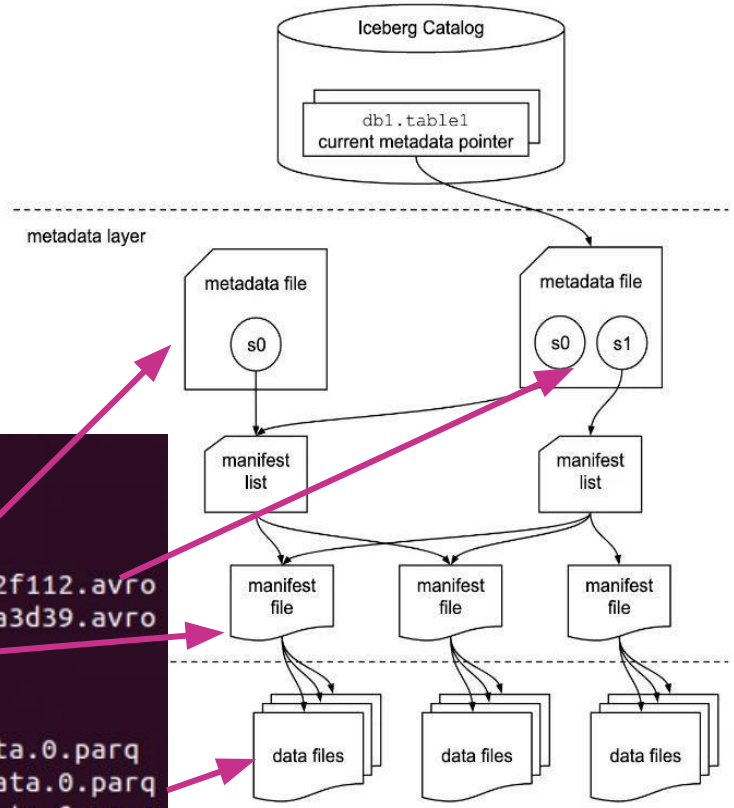
- Popular **table format**
- Defines how to:
  - Organize table data and metadata
  - Interact with meta/data -> Spec
- Table metadata on storage
- Famous features:
  - Flexible partitioning (transforms)
  - Partition/schema evolution
  - Time travel
  - Branching and tagging
  - Row-level modifications
- **Library/API**
  - Clients can interact with tables
- **Catalogs**
  - HMS, Glue, JDBC Rest (Polaris)



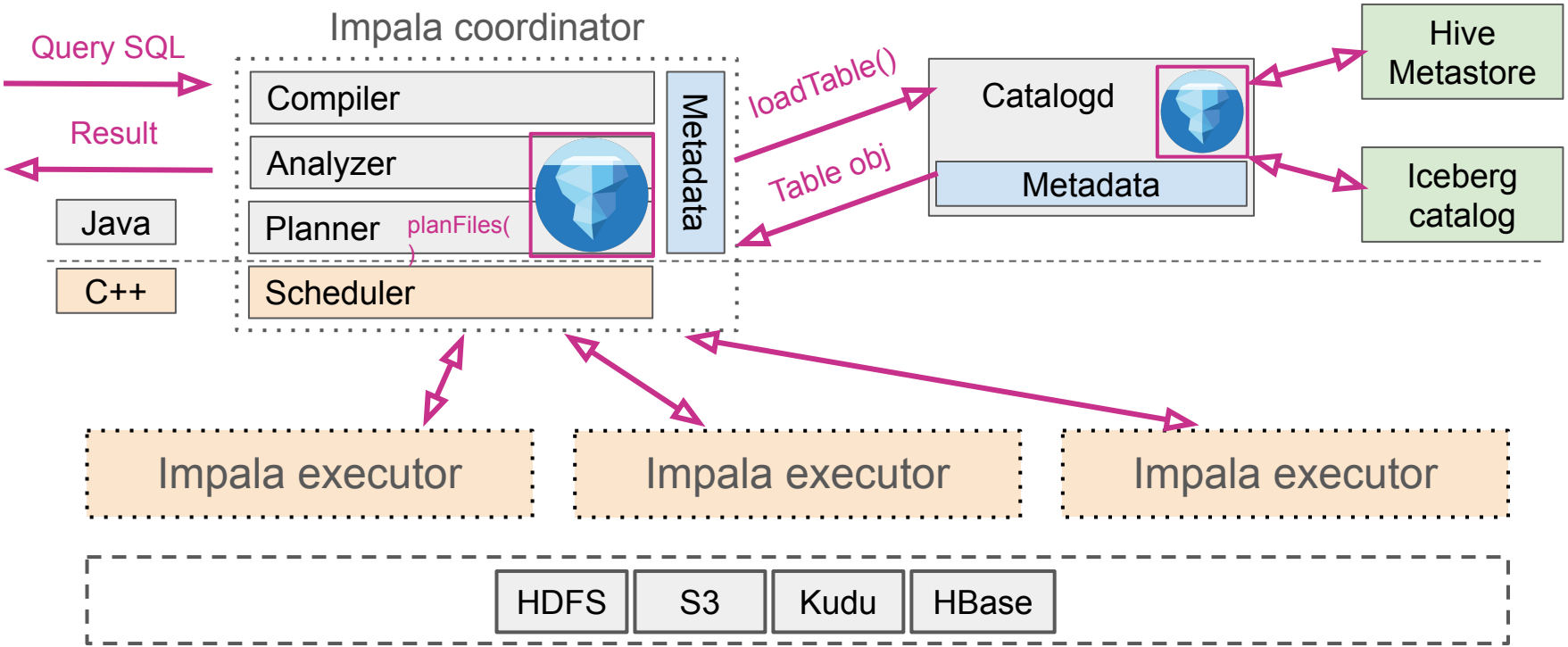
1. Introduction - Iceberg

- CREATE TABLE tbl (i int, s string)
  PARTITIONED BY SPEC (truncate(3, s))
  STORED AS ICEBERG
  TBLPROPERTIES ('format-version'='2');
- INSERT INTO tbl VALUES (1, "**abc**d"), (2, "**xyz**1");
- INSERT INTO tbl VALUES (3, "**abc**xyz");

```
tbl/metadata/
  00000-7e01eda3-380a-4d83-9416-050cec97ef81.metadata.json
  00001-212fafed-3bf0-4f91-beb8-835969c4b13c.metadata.json
  00002-7f081e0a-7a0f-4aa8-aa3e-99f35e97658b.metadata.json
  snap-3990482029540480076-1-1f831dc7-16bb-4490-8354-594717d2f112.avro
  snap-8137342376748057061-1-628a9e5a-c146-4a93-96d7-cd3a546a3d39.avro
  1f831dc7-16bb-4490-8354-594717d2f112-m0.avro
  628a9e5a-c146-4a93-96d7-cd3a546a3d39-m0.avro
tbl/data/
  s_trunc=abc/1c470c37d3f7cf65-c8fbfa3800000000_558329292_data.0.parq
  s_trunc=abc/e443b0000d3885ce-57be348300000000_2116373773_data.0.parq
  s_trunc=xyz/1c470c37d3f7cf65-c8fbfa3800000000_1332670711_data.0.parq
```



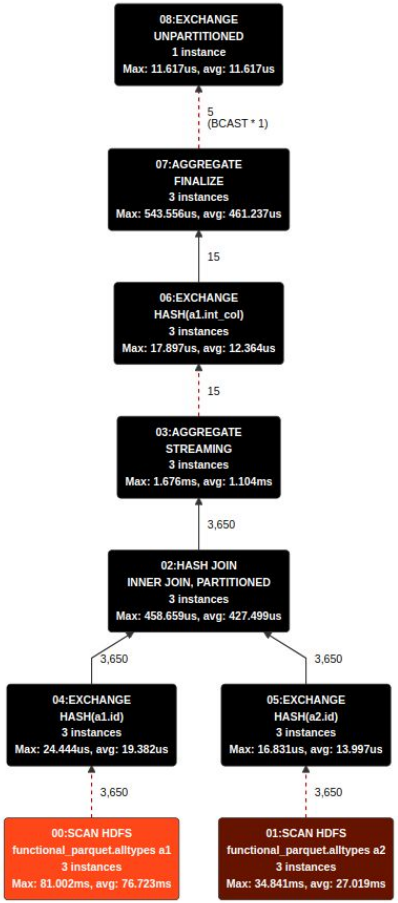
1. Introduction - Impala



1. Introduction - Impala

- Query plan example

```
SELECT count(1), avg(a1.int_col)
FROM
    functional_parquet.alltypes a1,
    functional_parquet.alltypes a2
WHERE
    a1.id = a2.id AND
    a1.id % 2 = 0
GROUP BY a1.int_col;
```



Contents

2. Row-level deletes - Concepts

DELETE FROM tbl WHERE id = 15;

Merge-on-read

- Tracking deleted rows in a separate “delete file”
- Good for frequent, small modifications
- Low write amplification
- High read amplification
- Table maintenance is a MUST

Copy-on-write

- Replaces old data files with rewritten data files
- Useful for infrequent, large modifications
- High write amplification
- No read amplification

2. Row-level deletes - Concepts

DELETE FROM tbl WHERE id = 15;

Positional deletes

- File_path + position
- Slower writes
- Better perf. to read

‘path1/abc.parquet’	13
‘path1/abc.parquet’	1234
‘path2/xyz.parquet’	1

Equality deletes

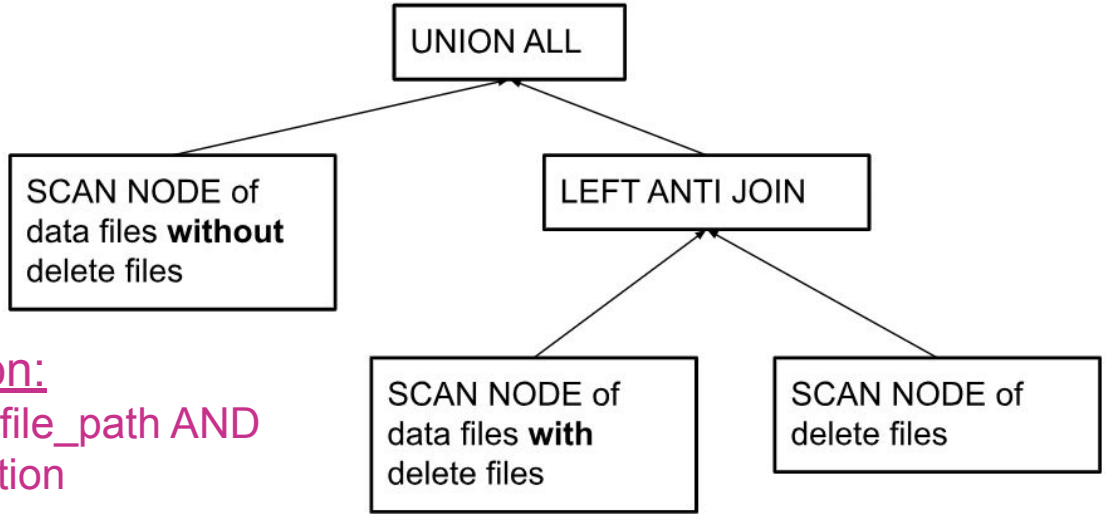
- Schema depends on ‘identifier-field-ids’
- Cheap to write
- Inefficient to read

ID
15

ID_col1	ID_col2	ID_col3
42	“string”	07.10.2024

2. Row-level deletes - Implementation

- Read data files w/o deletes
- Read delete files
- Read data file w/ deletes but add **virtual columns**



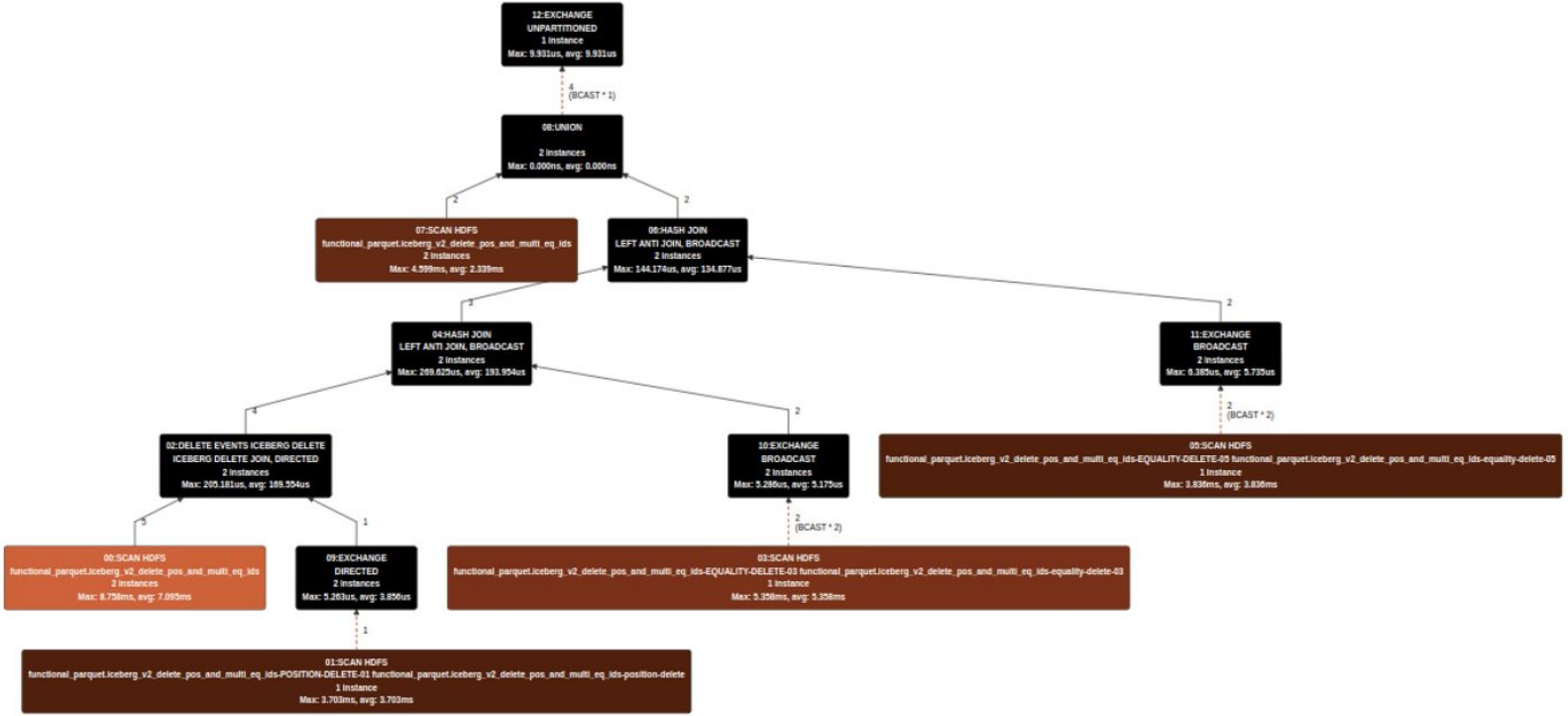
Position delete Anti-Join condition:

tbl.**INPUT__FILE__PATH** = DEL-tbl.file_path AND
tbl.**FILE__POSITION** = DEL-tbl.position

Equality delete Anti-Join condition:

tbl.ID_col1 IS NOT DISTINCT FROM DEL-tbl.ID_col1 AND … AND
tbl.DATA__SEQUENCE__NUMBER < DEL-tbl.DATA__SEQUENCE__NUMBER

2. Row-level deletes - Implementation



2. Row-level deletes - Pos delete read performance - 1

Details for performance measurement:

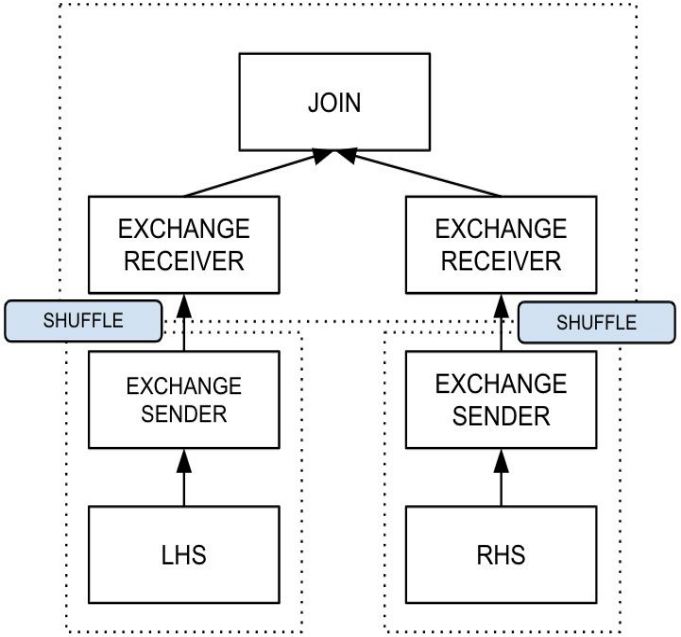
- Table with 8.64bn rows. ~10% deleted by position delete files.
- Query: SELECT count(1) ran approx **21 sec**

Operator	#Hosts	#Inst	Avg Time	Max Time	#Rows	Detail
F03:ROOT	1	1	0.000ns	0.000ns		
07:AGGREGATE	1	1	0.000ns	0.000ns	1	FINALIZE
06:EXCHANGE	1	1	0.000ns	0.000ns	480	UNPARTITIONED
F02:EXCHANGE SENDER	40	480	108.333us	4.000ms		
03:AGGREGATE	40	480	44.808ms	136.000ms	480	
02:DELETE EVENTS ICEBERG DELETE	40	480	715.567ms	2s347ms	7.81B	ICEBERG DELETE JOIN, PARTITIONED
|--F04:JOIN BUILD	40	480	226.358ms	739.997ms		
| 05:EXCHANGE	40	480	31.233ms	103.999ms	825.05M	HASH(<tbl_name>-delete.file_path)
| F01:EXCHANGE SENDER	40	478		3s716ms		
| 01:SCAN S3	40	478	365.866ms	1s036ms	825.05M	<tbl_name>-position-delete
04:EXCHANGE	40	480	339.825ms	1s047ms	8.64B	HASH(<tbl_name>.input__file__name)
F00:EXCHANGE SENDER	40	480		6s784ms		
00:SCAN S3	40	480	382.300ms	560.003ms	8.64B	<tbl_name>

2. Row-level deletes - Pos delete read performance - 1

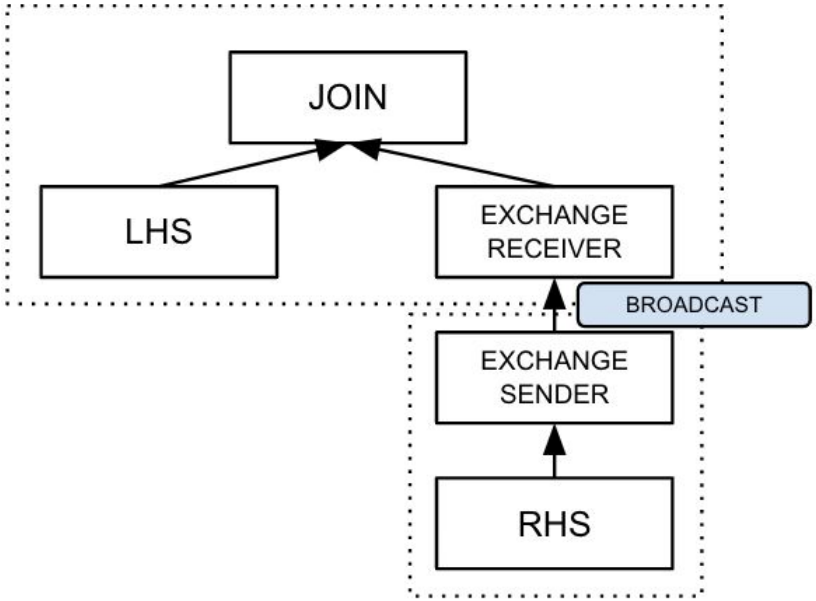
Partitioned

Cost:
#rows_left_scan + rows_right_scan



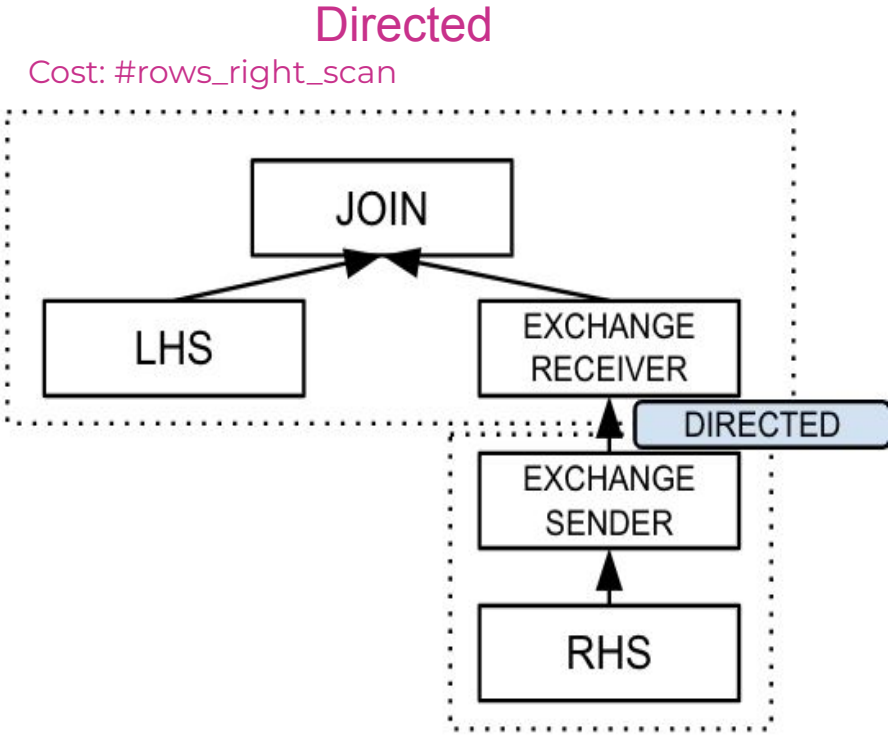
Broadcast

Cost: #rows_right_scan * #JOIN_nodes



2. Row-level deletes - Pos delete read performance - 1

- **DIRECTED** distribution mode:
  Use '**file_path**' in delete file to route rows
- No need to send 'left' rows
- No need to broadcast 'right' rows
- Cost: #rows_right_scan
- **Reduced** query runtime by **42%**



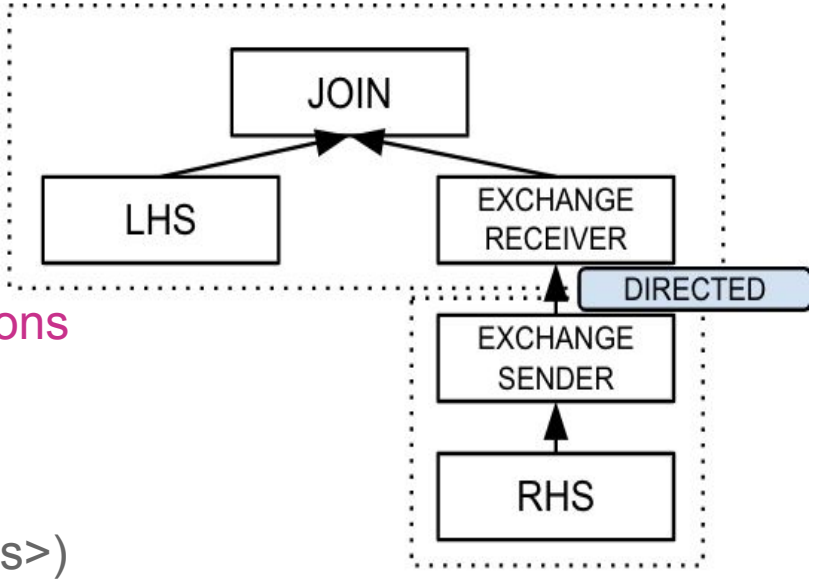
2. Row-level deletes - Pos delete read performance - 2

Another measurement:

- Table with 1 trillion rows. ~68.5bn rows deleted by position delete files.
- Query: SELECT count(1) ran approx **7m15s**

Operator	#Hosts	#Inst	Avg Time	Max Time
04:UNION	40	480	321.775ms	440.002ms
|--02:DELETE EVENTS ICEBERG DELETE	40	480	15s997ms	18s288ms
| |--F06:JOIN BUILD	40	40	4m15s	4m46s
| | 07:EXCHANGE	40	40	**47s784ms**	**58s928ms**
| | F02:EXCHANGE SENDER	40	480	**36s512ms**	**57s008ms**
| | 01:SCAN S3	40	480	12s161ms	20s308ms
| 00:SCAN S3	40	480	20s370ms	23s792ms
03:SCAN S3	40	480	42s696ms	46s015ms

2. Row-level deletes - Pos delete read performance - 2

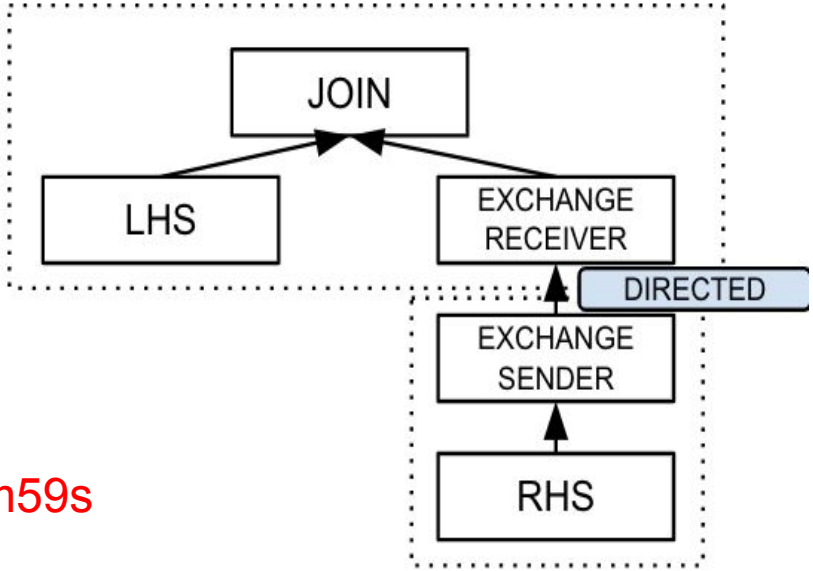


JOIN BUILD: 4m46s

- We build hashmap:
  ("file path" -> vector<positions>)
- Many re-allocation when adding positions to the vector

- Instead build hashmap:
  ("file_path" -> RoaringBitmap<positions>)
- Join build from 4m46s -> 1m49s
- Query runtime: 7m15s -> 4m59s

2. Row-level deletes - Pos delete read performance - 3



EXCHANGE SENDER: ~1m

- File paths being sent redundantly
- Pos dels ordered by file path.
- Remove redundancy

- Exchange sender from ~1m -> ~21.5s
- Query runtime: 7m15s -> 4m59s -> 3m59s

StringVal	int64	“long_file_path.parq”	StringVal	int64	“long_file_path.parq”	StringVal	int64	“long_file_path.parq”

“long_file_path”	StringVal	int64	StringVal	int64	StringVal	int64

Contents

3. Metadata table queries - Implementation

Iceberg API to query metadata tables:

data_files
delete_files
entries
files
manifests

history
metadata_log_entrie
partitions
position_deletes
refs
snapshots

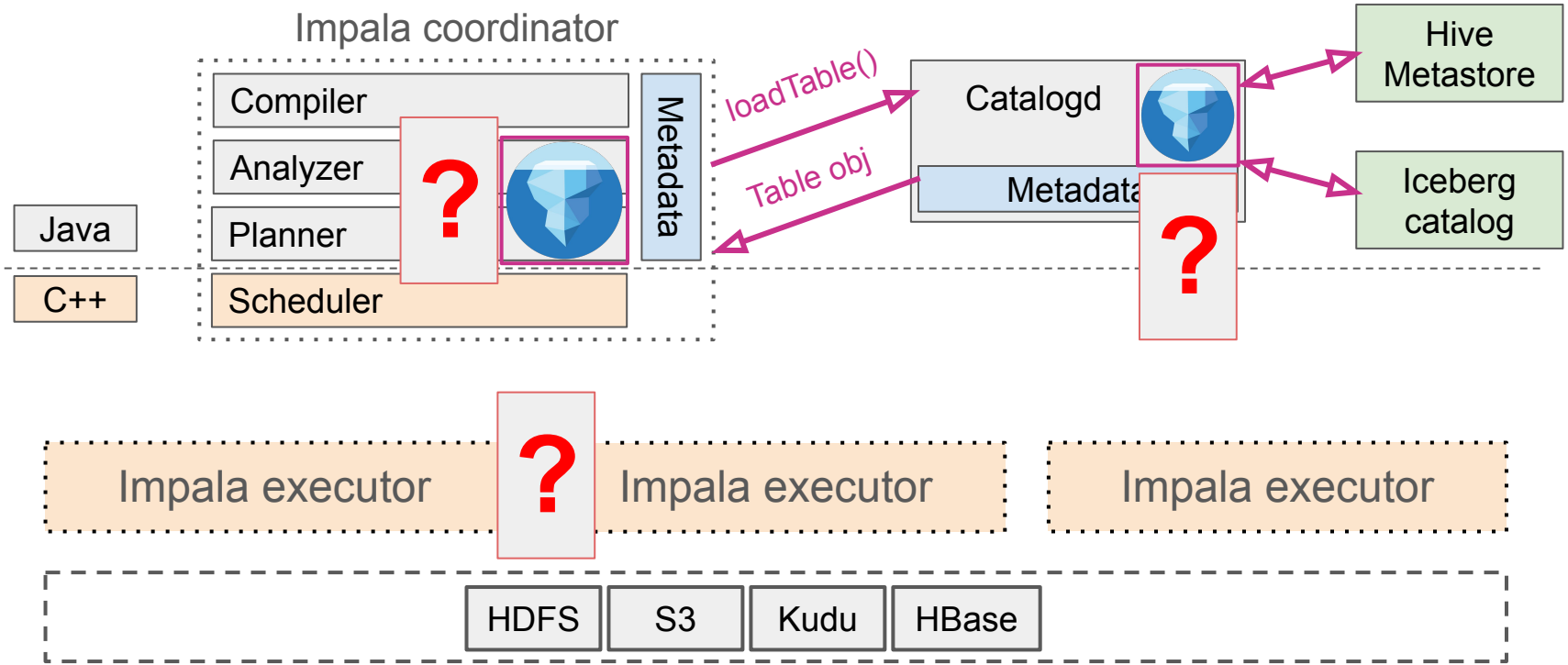
all_data_files
all_delete_files
all_entries
all_fi
all_n

```
SELECT
    s.operation,
```

```
Table metaTbl = MetadataTableUtils.createMetadataTableInstance(
    tbl, /* An Iceberg table object */
    MetadataTableType.PARTITIONS);
for (FileScanTask task : metaTbl.newScan().planFiles()) {
  for (StructLike row : ((DataScan)task).rows()) {
    // Get fields from 'row'
  }
}
```

```
ry h
nots s
```

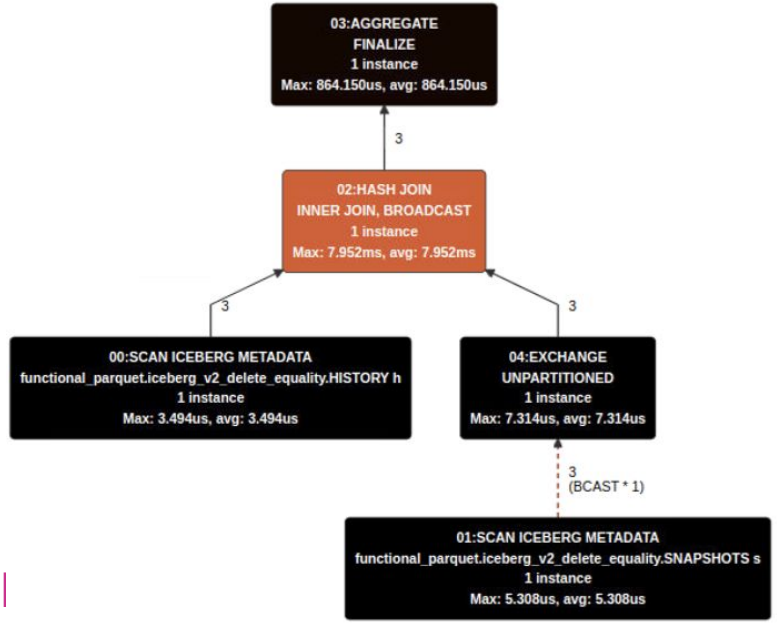
3. Metadata table queries - Implementation



3. Metadata table queries - Implementation

Thought process:

- Implement metadata SCAN in Executor?
  - SCAN would fit into Impala's architecture (plan tree).
  - Would need C++ Iceberg API or implement reads for ourselves
- Answer metadata SCAN in coordinator?
  - It's Java, simple to implement
  - Can't do 'regular query' functionality like joining, aggregating, etc.
- Still, should do the SCAN as part of the pl tree



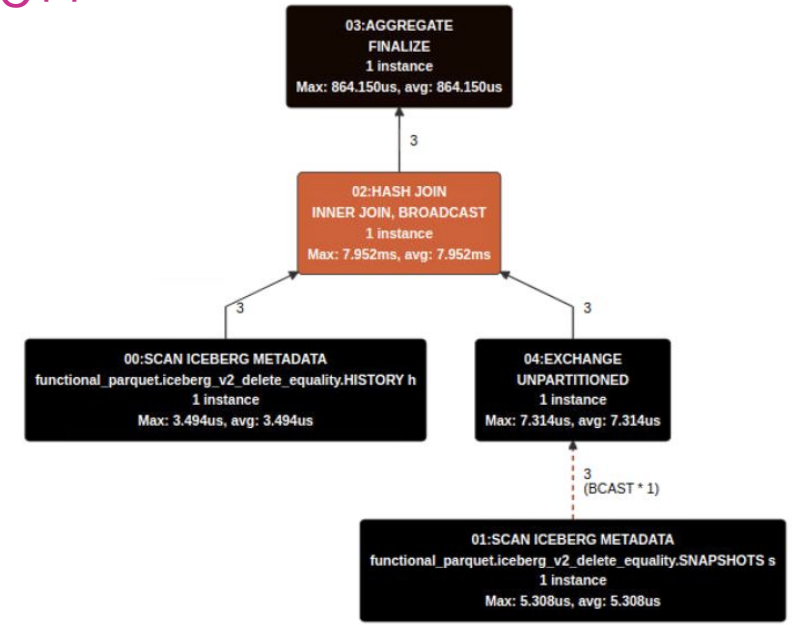
3. Metadata table queries - Implementation

Solution:

- SCAN ICEBERG METADATA node on C++ side
- Metadata Scanner on Java side
- JNI call from C++ to Java to get rows

Trade-offs and difficulties:

- Metadata SCANs are coordinator only
  - There is C++ and Java too
- Beware! GC vs access from C++
- Type conversion from Java to C++
- Extra steps to populate ‘RowBatch‘
- Code readability
- Performance?



Contents

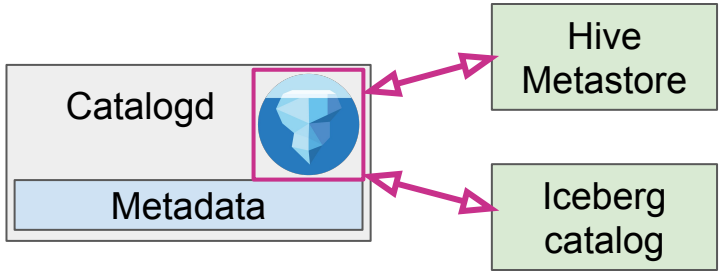
4. Catalogs - Currently

Iceberg catalogs:

- HiveCatalog, HadoopCatalog, JdbcCatalog, NessieCatalog, RestCatalog, GlueCatalog, SnowflakeCatalog, etc.

Catalogs supported by Impala:

- HiveCatalog and HadoopCatalog (non-prod)

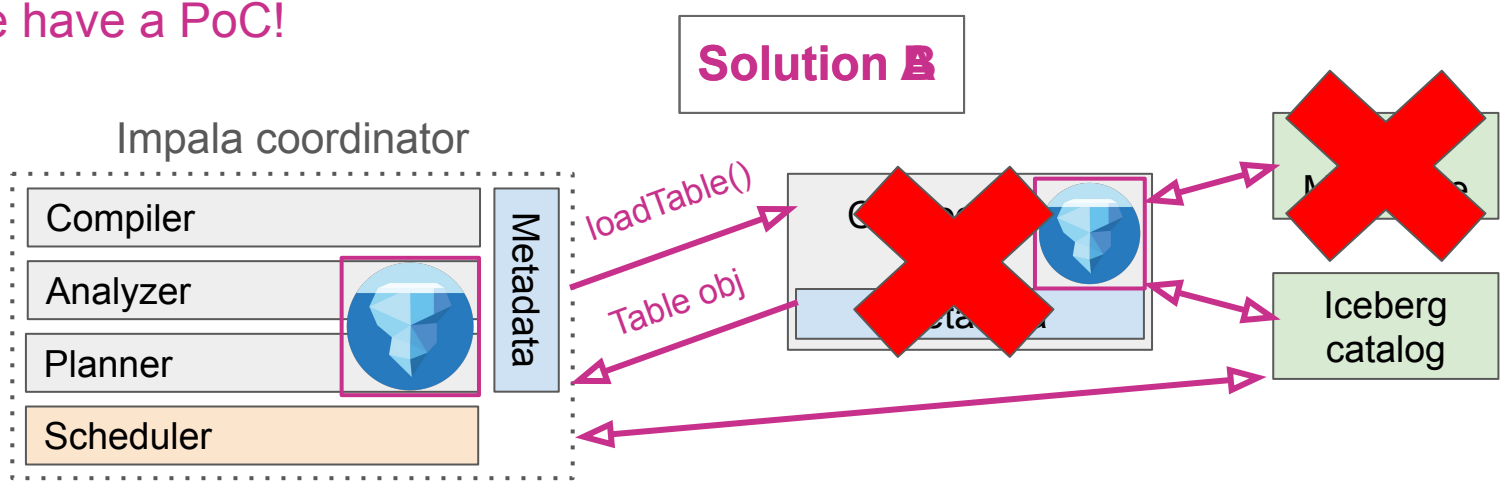


Other limitations:

- Heavy HMS dependency
- Full table name: DB.TBL instead of CATALOG.DB.TBL
- No flexibility for configuration

4. Catalogs - Future plans

- More catalog types: RESTCatalog!
- Catalog abstraction on top of DB.TBL
- More flexible creation + configuration
- Reduce HMS dependency
- We have a PoC!



Contents

+1 Iceberg V3 positional deletes

Proposal for new Positional delete design

V2 Positional deletes

- File_path + position
- Single delete file for multiple data files
- New deletes for writes

‘path1/abc.parquet’	13
‘path1/abc.parquet’	1234
‘path2/xyz.parquet’	1

V3 Positional deletes

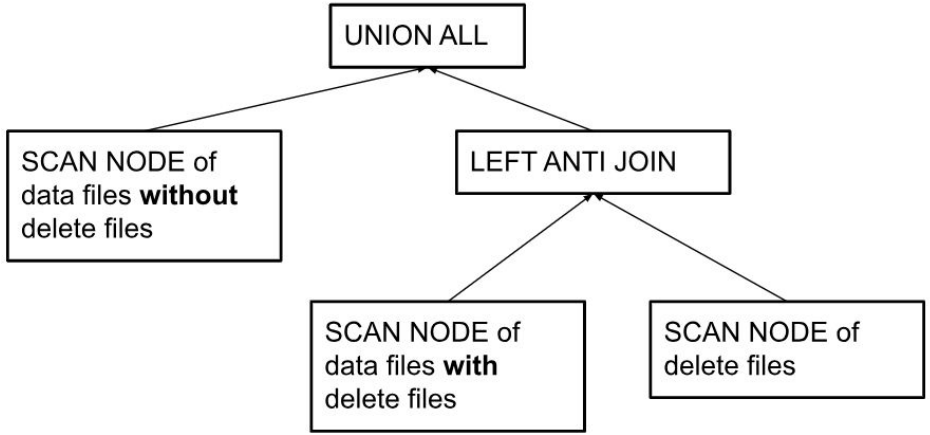
- Delete vector as a RoaringBitmap
- One delete vector for one data file
- Multiple bitmaps in a Puffin file
- File path + offset + length stored in Iceberg metadata
- Merge bitmaps for writes

[Puffin header][bitmap1]...[bitmapN][Puffin footer]

+1 Iceberg V3 positional deletes

Exploring opportunities

1) Delete SCAN node to read bitmaps and return as if V2
2) Delete SCAN to read bitmaps and return bitmaps
3) No Delete SCAN node, JOIN BUILD to read bitmaps



Difficulties

- Need a C++ Puffin reader and writer
- Merge bitmaps before writing
- Cross language compatibility?

Questions?